

Persistence and Monotony Properties of Argumentation Semantics

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Overview

- 1 Introduction
- 2 Preliminaries
- 3 Addition Persistence
- 4 Removal Persistence
- 5 Skeptical Monotony
- 6 Some Remarks About the Stable Semantics
- 7 Conclusions and Future Research

- Argumentation is a dynamic process,
 - but an argumentation framework represents only a static snapshot.
- To consider dynamics, we need to consider AFs that change,
 - e.g. addition of new arguments and attacks.
- **How do semantics for argumentation behave when the AF changes?**
- 2009: Overview for grounded semantics (Boella, Kaci, vdT)
- Afterwards: a lot of work on dynamics, abduction, counterfactuals, revision, strong equivalence, etc etc
 - But no similar overview of properties for multiple extensions?
 - NB: If you know some other overview, let me know!

- The main question we address is:

When is a point of view on argument acceptance robust w.r.t. addition/removal of attacks?

- We break this down by considering the following three properties:
 - ① **XY addition persistence:** A σ labelling of an AF F in which x is labelled X and y is labelled Y is still a σ labelling of F after adding an attack from x to y .
 - ② **XY removal persistence:** A σ labelling of an AF F in which x is labelled X and y is labelled Y is still a σ labelling of F if removing the attack from x to y .
 - ③ **XY skeptical monotony:** If in all σ labellings of an AF F , x is labelled X and y is labelled Y , then adding an attack from x to y does not lead to new σ labellings.

(Alternative: only consider arguments labeled in.)

Definition

Let \mathcal{U} be a set whose elements are called *arguments*. An *argumentation framework* is a pair $F = (A, \rightsquigarrow)$ where A is a finite subset of \mathcal{U} and $\rightsquigarrow \subseteq A \times A$ is a relation called the *attack relation*. We denote by \mathcal{F} the set of all argumentation frameworks.

Definition

A *labelling* of an argumentation framework (A, \rightsquigarrow) is a function $L : A \rightarrow \{\mathbf{I}, \mathbf{O}, \mathbf{U}\}$. Given a label $l \in \{\mathbf{I}, \mathbf{O}, \mathbf{U}\}$ we define $L^{-1}(l)$ as $\{x \in A \mid L(x) = l\}$. Given an argumentation framework F , we let $\mathcal{L}(F)$ denote the set of all labellings of F .

Definition

Let $F = (A, \rightsquigarrow)$ be an argumentation framework. A labelling $L \in \mathcal{L}(F)$ is *complete* if and only if:

- 1 for all $x \in A$, $L(x) = \mathbf{I}$ iff for all y s.t. $y \rightsquigarrow x$, $L(y) = \mathbf{O}$.
- 2 for all $x \in A$, $L(x) = \mathbf{O}$ iff for some y s.t. $y \rightsquigarrow x$, $L(y) = \mathbf{I}$.

The complete, grounded, preferred, semi-stable and stable semantics are defined as follows.

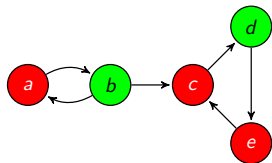
Definition

Let $F = (A, \rightsquigarrow)$ be an argumentation framework.

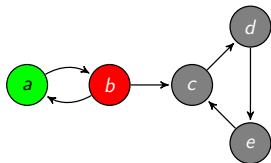
- $\mathcal{L}_{co}(F) = \{L \in \mathcal{L}(F) \mid L \text{ is a complete labelling of } F\}$
- $\mathcal{L}_{gr}(F) = \{L \in \mathcal{L}_{co}(F) \mid \nexists L' \in \mathcal{L}_{co}(F) \text{ s.t. } L'^{-1}(\mathbf{I}) \subset L^{-1}(\mathbf{I})\}$
- $\mathcal{L}_{pr}(F) = \{L \in \mathcal{L}_{co}(F) \mid \nexists L' \in \mathcal{L}_{co}(F) \text{ s.t. } L^{-1}(\mathbf{I}) \subset L'^{-1}(\mathbf{I})\}$
- $\mathcal{L}_{ss}(F) = \{L \in \mathcal{L}_{co}(F) \mid \nexists L' \in \mathcal{L}_{co}(F) \text{ s.t. } L'^{-1}(\mathbf{U}) \subset L^{-1}(\mathbf{U})\}$
- $\mathcal{L}_{st}(F) = \{L \in \mathcal{L}_{co}(F) \mid L^{-1}(\mathbf{U}) = \emptyset\}$

An example

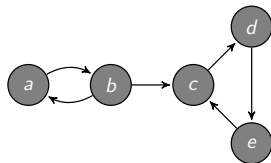
We use green, red and grey to depict in, out and undecided arguments.



: L_1



: L_2



: L_3

	Complete	Grounded	Preferred	Semi-Stable	Stable
L_1	✓		✓	✓	✓
L_2	✓		✓		
L_3	✓	✓			

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Definition (XY Addition Persistence)

A semantics σ satisfies XY addition persistence iff every σ labelling of an AF F in which x is labelled X and y is labelled Y is still a σ labelling of F after adding an attack from x to y .

Note: only grounded, complete, preferred and semi-stable semantics.

We discuss stable semantics later (roughly: trivial or same as preferred).

Addition Persistence Properties

Some addition persistence properties fail simply because they introduce a conflict. The following examples apply to all semantics we consider.

Failure of **II**-addition persistence



Failure of **IU**-addition persistence



Failure of **UI**-addition persistence



Addition Persistence Properties

As we just saw, some properties fail:

- **II**-addition persistence
 - **IU**-addition persistence
 - **UI**-addition persistence
- } These cases fail because they introduce a conflict.

Other properties reflect reasonable principles:

- **OO**-addition persistence
 - **OU**-addition persistence
 - **OI**-addition persistence
 - **IO**-addition persistence
 - **UO**-addition persistence
 - **UU**-addition persistence
- } In these cases the added attack doesn't introduce a conflict, and doesn't invalidate the justification of the attacked argument's label.

Are these properties satisfied by the semantics we consider?

Addition Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	✓	-	-

Preferred:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	✓	-	-

Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

Addition Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	✓	-	-

Preferred:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	✓	-	-

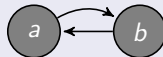
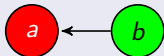
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

Addition Persistence Properties

Failure of **OI** addition persistence under grounded semantics.

After adding an attack from a to b there is a new grounded labelling:



Addition Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	✓	-	-

Preferred:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	✓	-	-

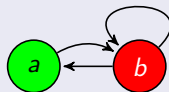
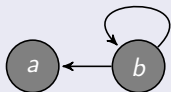
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

Addition Persistence Properties

Failure of **UU** addition persistence under the preferred semantics.

After adding an attack from a to b there is a new preferred labelling:



Addition Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	✓	-	-

Preferred:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	✓	-	-

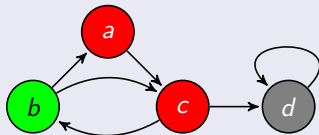
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

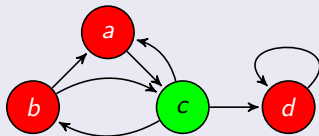
Addition Persistence Properties

Failure of \mathbf{OO} addition persistence under the semi-stable semantics

The following AF has one semi-stable labelling:



The arguments a and c are both \mathbf{O} . But if we add an attack $c \rightsquigarrow a$ we get a different semi-stable labelling:



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Definition (XY Removal Persistence)

A semantics σ satisfies XY removal persistence iff: every σ labelling of an AF F in which x is labelled X and y is labelled Y is still a σ labelling of F after removing the attack from x to y .

Note: again, we look only at the grounded, complete, preferred and semi-stable semantics. We discuss the stable semantics later.

Removal Persistence Properties

Some removal persistence properties fail because the removal of an attack from x to y may make the label of y unjustified. The following examples apply for all semantics we consider.

Failure of **IO**-removal persistence:



Failure of **UU**-removal persistence:



Removal Persistence Properties

As we just saw, some removal persistence properties fail:

- **IO**-removal persistence
 - **UU**-removal persistence
- } These cases fail because they invalidate the justification of the attacked argument's label.

Others are trivially satisfied:

- **II**-removal persistence
 - **UI**-removal persistence
 - **IU**-removal persistence
- } Combinations of labels that never occur.

The remaining properties reflect reasonable principles:

- **OO**-removal persistence
 - **OU**-removal persistence
 - **OI**-removal persistence
 - **UO**-removal persistence
- } In these cases the removal does not invalidate the attacked argument's label.

Are these properties satisfied by the semantics we consider?

Removal Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Complete:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Preferred:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	✓
	I	-	✓	✓

Removal Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Complete:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Preferred:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

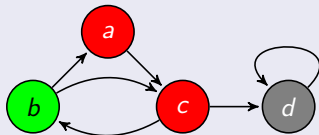
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	✓
	I	-	✓	✓

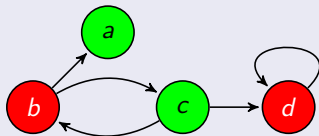
Removal Persistence Properties

Failure of **OO** removal persistence under the semi-stable semantics

The following AF has one semi-stable labelling:



The arguments a and c are both **O**. But if we remove the attack $a \rightsquigarrow c$ we get a different semi-stable labelling:



Removal Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Complete:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Preferred:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

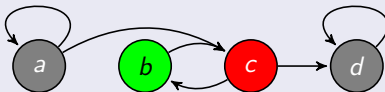
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	✓
	I	-	✓	✓

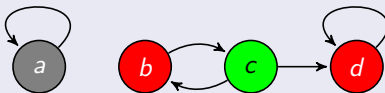
Removal Persistence Properties

Failure of **UO** removal persistence under the semi-stable semantics

The following AF has one semi-stable labelling:



The arguments a and c are labelled **U** and **O**. If we remove the attack $a \rightsquigarrow c$ we get a different semi-stable labelling:



Removal Persistence Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Complete:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

Preferred:

		X		
		O	U	I
Y	O	✓	✓	-
	U	✓	-	✓
	I	✓	✓	✓

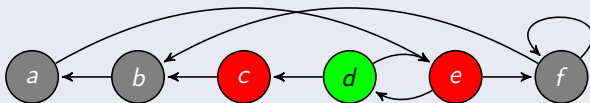
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	✓
	I	-	✓	✓

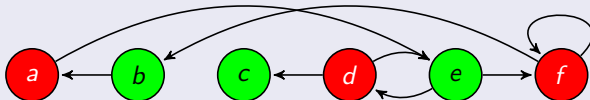
Removal Persistence Properties

Failure of **OU** removal persistence under the semi-stable semantics

The following AF has one semi-stable labelling:



The arguments *c* and *b* are labelled **O** and **U**. If we remove the attack $c \rightsquigarrow b$ we get a different semi-stable labelling:



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Skeptical Monotony Properties

Intuition

Suppose that x and y are labelled X and Y in *all* σ labellings of F . The XY addition persistence property then implies that all σ labellings of F are still σ labellings after adding the attack $x \rightsquigarrow y$. In other words, *no σ labelling gets destroyed*. But is it also the case that *no new labellings are created*? This is the property that we consider here.

Definition (Skeptical XY monotony)

A semantics σ satisfies skeptical XY monotony iff: If in all σ labellings of an AF F , x is labelled X and y is labelled Y , then **adding an attack from x to y does not lead to new σ labellings**.

Note: again, we look only at the grounded, complete, preferred and semi-stable semantics. We discuss the stable semantics later.

Skeptical Monotony Properties

Some skeptical monotony properties fail simply because they introduce a conflict and thus lead to new labellings. The following examples apply to all semantics we consider.

Failure of **II** skeptical monotony



Failure of **IU** skeptical monotony



Failure of **UI** skeptical monotony



Skeptical Monotony Properties

The previous examples show that the following properties fail:

- **II** skeptical monotony
 - **IU** skeptical monotony
 - **UI** skeptical monotony
- } These cases fail because they introduce a conflict and thus produce new labellings.

What about the other properties?

- **OO** skeptical monotony
- **OU** skeptical monotony
- **OI** skeptical monotony
- **IO** skeptical monotony
- **UO** skeptical monotony
- **UU** skeptical monotony

Addition Persistence Properties (reminder)

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	✓	-	-

Preferred:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	✓	-	-

Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

Skeptical Monotony Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	-	-	-

Preferred:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

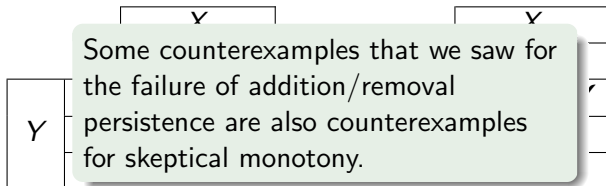
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

Skeptical Monotony Properties

Grounded:

Complete:



Preferred:

Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

Skeptical Monotony Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	-	-	-

Preferred:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

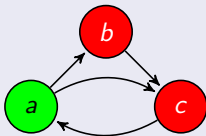
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

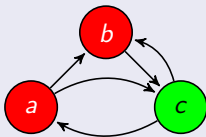
Skeptical Monotony Properties

Failure of **OO** skeptical monotony under the preferred semantics

The following AF has one preferred labelling:



The arguments b and c are both labelled **O**. If we add an attack $c \rightsquigarrow b$ then we obtain a new preferred labelling:



Skeptical Monotony Properties

Grounded:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	✓	-
	I	-	-	-

Complete:

		X		
		O	U	I
Y	O	✓	✓	✓
	U	✓	-	-
	I	-	-	-

Preferred:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

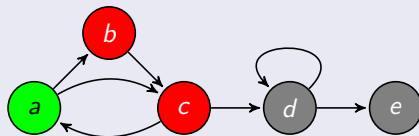
Semi-Stable:

		X		
		O	U	I
Y	O	-	-	-
	U	-	-	-
	I	-	-	-

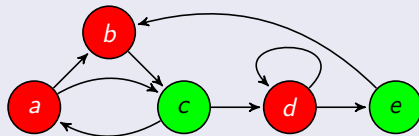
Skeptical Monotony Properties

Failure of **UO** skeptical monotony under the preferred semantics

The following AF has one preferred labelling:



The arguments *e* and *b* are labelled **U** and **O**. If we add an attack $e \rightsquigarrow b$ then we obtain a new preferred labelling:



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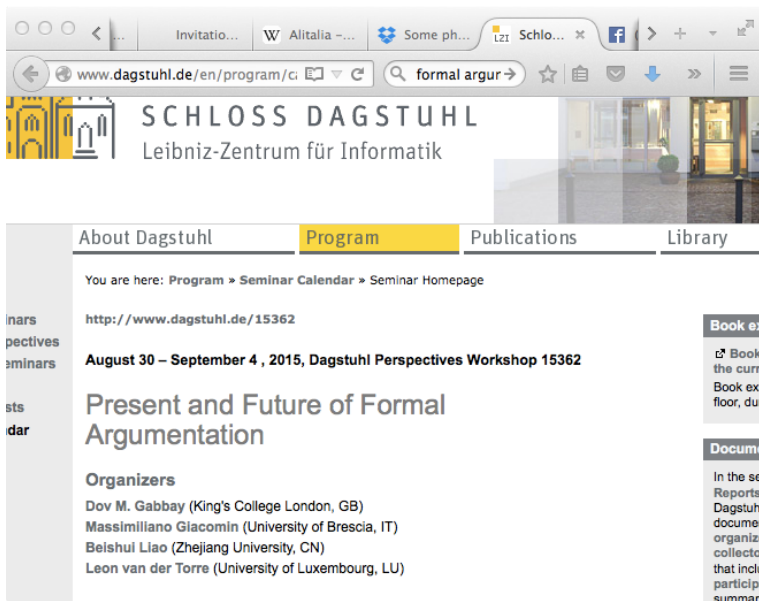
Some Remarks About the Stable Semantics

- All properties that we considered that involve **U** labelled arguments are trivially satisfied under the stable semantics, because no argument is ever labelled **U** under the stable semantics.
- For all properties that do not involve **U** labelled arguments, the results for the stable semantics are *exactly the same as for the preferred semantics*.

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- We studied the behaviour of semantics for argumentation when the AF changes.
- Some insights we gained are that some properties that appear intuitive are not satisfied under all semantics
 - Failure of e.g. **OO** skeptical monotony under the preferred semantics.
 - Failure of *many* properties under the semi-stable semantics.
- The complete semantics is the best behaved semantics w.r.t. the properties we discussed.


- Study skeptical removal monotony.
- Study weaker versions of the properties considered here.
- Study semantics not considered here.
- Study relationships with other work on the behaviour of semantics:
 - Strong equivalence.
 - Input/output behaviour.
 - The directionality principle.
- Study the consequences of our results in the setting of strategic argumentation and revision, counterfactuals and abduction in argumentation.



The screenshot shows a web browser window with the URL www.dagstuhl.de/en/program/ci. The page header features the logo of Schloss Dagstuhl Leibniz-Zentrum für Informatik and a photograph of the building entrance. A navigation menu includes 'About Dagstuhl', 'Program' (highlighted), 'Publications', and 'Library'. The breadcrumb trail reads 'You are here: Program > Seminar Calendar > Seminar Homepage'. The URL <http://www.dagstuhl.de/15362> is displayed. The main heading is 'August 30 – September 4, 2015, Dagstuhl Perspectives Workshop 15362' followed by 'Present and Future of Formal Argumentation'. The organizers listed are Dov M. Gabbay (King's College London, GB), Massimiliano Giacomin (University of Brescia, IT), Beishui Liao (Zhejiang University, CN), and Leon van der Torre (University of Luxembourg, LU). On the right side, there are partial views of 'Book ex...' and 'Docume...' sections.

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August 30 – September 4, 2015, Dagstuhl Perspectives Workshop 15362

Present and Future of Formal Argumentation

Organizers

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