

Exploring Relations between Answer Set Programs

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Relations between Answer Set Programs

⌘ **LP**: the class of all logic programs

⌘ Relation $\odot \subseteq \mathbf{LP} \times \mathbf{LP}$

⌘ $P1, P2 \in \mathbf{LP}$: logic programs

⌘ $P1 \odot P2$: comparison, correspondence, entailment, etc.

⌘ **Equivalence**/Non-equivalence

⌘ **Inclusion**/Non-inclusion/Disjoint

⌘ **Generality**/Specificity and Abstraction/Refinement

⌘ **Preferred**/Non-preferred and Higher/Lower Priority

⌘ **Strength**/Weakness and Better/Less Valued

⌘ More/Less **Informativeness** or Supportedness

Ordering Logic Programs

- ⌘ Equivalence of logic programs
 - ⌘ optimization, debugging, simplification, verification
 - ⌘ ordinary (weak)/strong/uniform/relativized/projected
- ⌘ Direction to refine/abstract programs (minimal revision)
 - ⌘ Inductive Logic Programming (ILP), Ontology
- ⌘ Synthesis of a common generalized/specialized theory from different sources of information [Sakama & Inoue, 2006,2007,2008]
 - ⌘ Multi-Agent Systems (MAS)

Comparing Nonmonotonic Programs

⌘ Example:

$P1 = \{ p \leftarrow \textit{not} q \}.$

$P2 = \{ p \leftarrow \textit{not} q, \quad q \leftarrow \textit{not} p \}.$

P1 has the single answer set: $\{p\}$

P2 has two answer sets: $\{p\}, \{q\}$

⌘ P1 is **more informative** than P2 in the sense that P1 has the **skeptical** consequences $\{p\}$ which includes $\{\}$.

⌘ P2 is **more informative** than P1 in the sense that P2 has the **credulous** consequences $\{p,q\}$ which includes $\{p\}$.

Generality Relations [Inoue & Sakama, 2006]

- ⌘ **Smyth** ($\#$) and **Hoare** (\flat) orderings between sets of answer sets.
- ⌘ Both $\#$ - and \flat -generalities are defined in a way that **weakly equivalent programs belong to the same equivalence class** induced by these orderings.
- ⌘ Both **minimal upper** and **maximal lower bounds** can be defined for any pair of programs in these generality orderings.
- ⌘ $\#$ -general programs entail **more skeptical consequences**, while \flat - general programs entail **more credulous consequences**.
- ⌘ Both **strong $\#$ -** and **strong \flat -generalities** are defined in a way that **strongly equivalent programs belong to the same equivalence class** induced by these orderings.
- ⌘ The proposed orderings can be applied to the class of **extended disjunctive programs** (EDPs).

Extensions (in the paper)

- ⌘ General framework to compare semantic structures—comparing **composite sets**, which are sets of literal sets, under any **pre-order** between literal sets.
- ⌘ Application of the comparison framework to any class of logic programs including **nested programs**, programs with **aggregates**, and any **propositional formulas** as well as under other semantics.
- ⌘ Variations of generality relations — **strong/uniform/relativized equivalence/inclusion/generality**
 - ⌘ comparison of generality with respect to **contexts—robustness**
 - ⌘ satisfaction of classical **inductive generality**
 - ⌘ application to formalizing **abductive generality**

Ordering on Powersets

⌘ *pre-order* \leq : binary relation which is reflexive and transitive

⌘ *partial order* \leq : pre-order which is also anti-symmetric

⌘ $\langle D, \leq \rangle$: *pre-ordered set / poset*

⌘ $\mathbf{P}(D)$: the power set of D

⌘ The **Smyth order**: for $X, Y \in \mathbf{P}(D)$,

$$X \vDash^{\#} Y \text{ iff } \forall x \in X \exists y \in Y. y \leq x$$

⌘ The **Hoare order**: for $X, Y \in \mathbf{P}(D)$,

$$X \vDash^b Y \text{ iff } \forall y \in Y \exists x \in X. y \leq x$$

⌘ Both $\langle \mathbf{P}(D), \vDash^{\#} \rangle$ and $\langle \mathbf{P}(D), \vDash^b \rangle$ are pre-ordered sets.

Ordering over the sets of literal sets

- ⌘ Lit : the set of all ground literals in the language
- ⌘ **composite set** Σ : $\Sigma \in \mathbf{P}(\mathbf{P}(Lit))$
- ⌘ A composite set represents a semantic structure of a logic program P , e.g., the **answer sets (stable models)**, the **supported models**, the **possible models**, the **minimal models**, the **classical models**, the **preferred answer sets** of P .
- ⌘ The **Smyth order** $\Sigma1 \models^\# \Sigma2$ and the **Hoare order** $\Sigma1 \models^b \Sigma2$ can be defined for any two composite sets $\Sigma1, \Sigma2$.
- ⌘ If $\Sigma1$ and $\Sigma2$ represent the semantics of $P1$ and $P2$, the Smyth/Hoare ordering gives **generality relation** between them.

General Extended Disjunctive Programs (or simply Programs)

⌘ Rules:

$$L_1 ; \dots ; L_k ; \textit{not} L_{k+1} ; \dots ; \textit{not} L_l$$
$$\leftarrow L_{l+1}, \dots, L_m, \textit{not} L_{m+1}, \dots, \textit{not} L_n$$

- *not* : negation as failure, L_i : literal, $0 \leq k \leq l \leq m \leq n$
- **Extended disjunctive program (EDP):** $k=l$.
- **Extended logic program (ELP):** $k=l \leq 1$.
- **Positive disjunctive program (PDP):**
EDP & $\forall L_i$: atom & $m=n$.
- **Normal logic program (NLP):** ELP & $\forall L_i$: atom.

Ordering Logic Programs

⌘ **LP**: the class of all programs

⌘ $A(P), A(Q)$: the answer sets of $P, Q \in \mathbf{LP}$

● P is **more #-general than** Q :

$$P \models^{\#} Q \text{ iff } A(P) \models^{\#} A(Q)$$

● P is **more b-general than** Q :

$$P \models^b Q \text{ iff } A(P) \models^b A(Q)$$

⌘ Theorem:

(1) $P \models^{\#} Q$ and $Q \models^{\#} P$ iff $\min(A(P)) = \min(A(Q))$.

(2) $P \models^b Q$ and $Q \models^b P$ iff $\max(A(P)) = \max(A(Q))$.

In case of EDP, $P \models^{\#/b} Q$ and $Q \models^{\#/b} P$ iff $A(P) = A(Q)$ (i.e. $P \equiv Q$).

Ordering Logic Programs

⌘ Example:

$$P1 = \{ p \leftarrow \textit{not} \ q \}.$$

$$P2 = \{ p \leftarrow \textit{not} \ q, \quad q \leftarrow \textit{not} \ p \}.$$

$$P3 = \{ p; q \leftarrow \}.$$

$$P4 = \{ p \leftarrow \textit{not} \ \neg p, \quad q \leftarrow p \}.$$

$$A(P1) = \{\{p\}\}, \quad A(P2) = A(P3) = \{\{p\}, \{q\}\}, \quad A(P4) = \{\{p,q\}\}$$

- $P4 \not\models^{\#} P1 \not\models^{\#} P2$

- $P4 \models^b P2 \models^b P1$

- $P2 \not\models^{\#} P3 \not\models^{\#} P2, \quad P2 \models^b P3 \models^b P2 \quad (\text{i.e., } P2 \equiv P3)$

Skeptically/Credulously Entailed Literals in More/Less General Programs

⌘ $Skp(P)$: the set of skeptical consequences of P

⌘ $Crd(P)$: the set of credulous consequences of P

⌘ Theorem:

● If $P \models^{\#} Q$ then $Skp(P) \supseteq Skp(Q)$.

● If $P \models^b Q$ then $Crd(P) \supseteq Crd(Q)$.

Strong/Uniform/Relativized Generality

⌘ $P, Q \in \mathbf{LP}, C \subseteq \mathbf{LP}$

● P is **strongly/uniformly more #-general than** Q :

$P \underline{\triangleright}^{\#} Q$ iff $P \cup R \not\models^{\#} Q \cup R$ for any $R \in \mathbf{LP} / R \in \mathbf{P}(Lit)$.

● P is **strongly/uniformly more b-general than** Q :

$P \underline{\triangleright}^b Q$ iff $P \cup R \not\models^b Q \cup R$ for any $R \in \mathbf{LP} / R \in \mathbf{P}(Lit)$.

⌘ $P \underline{\triangleright}^{\#/b} Q$ implies $P \not\models^{\#/b} Q$.

⌘ **Relativized #/b**-generality requests the above for any $R \in C$

⌘ Strong generality is equivalent to **strong inclusion** [Eiter *et al.*, 2005] (for finite programs).

⌘ Uniform generality for PDPs reduces to **classical inductive generality**, i.e., logical entailment.

Abductive Generality

- ⌘ **Explainable generality** [Inoue & Sakama, 2008] can be characterized by **b-generality**. That is, the ability to explain more observations corresponds to derive more credulous consequences when abducibles are represented as **choice constructs**:

$L ; \textit{not } L \leftarrow$ for any abducible literal L .

- ⌘ **Explanatory generality** [Inoue & Sakama, 2008] can be characterized by **relativized b-generality**, in which the set of abducibles is set to be the **context** for comparison.
- ⌘ Taking minimal upper/maximal lower bounds of two abductive programs can be applied to coordination of multiple **abductive agents**.

Conclusion: Generality as Comparison between Programs

- ⌘ **Program Verification**: testing and debugging with respect to specification of a program (**equivalence**)
- ⌘ **Program Development**: testing if a refined theory is more detailed from the previous theory (**generality**)
- ⌘ **Problem Solving**: assessment of relative value of each theory/ontology (**generality, preference**)
- ⌘ **Machine Learning**: criteria to compose better theories (**generality, preference**)
- ⌘ **Multiagent Systems**: criteria to compare different information sources (**informativeness, preference**)