

# Formalizing Commitments Using Action Languages

Tran Cao Son<sup>1</sup>, Enrico Pontelli<sup>1</sup>,  
Chiaki Sakama<sup>2</sup>

<sup>1</sup>New Mexico State University

<sup>2</sup>Wakayama University

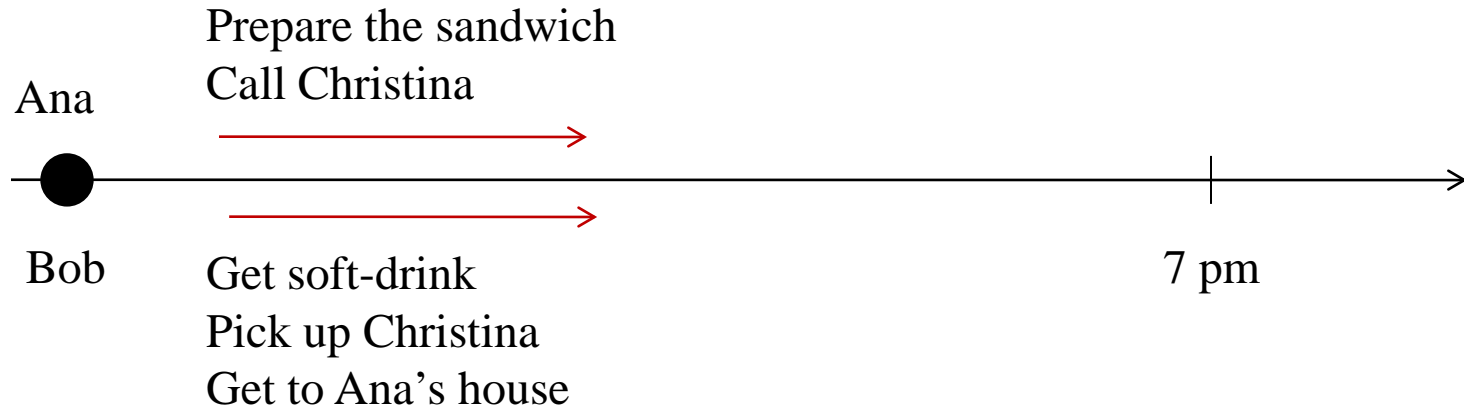
# Commitments

- Commitments are an integral part of agent societies.
  - Related to agents' behavior and capabilities
  - Often associated with time constraints in the future
  - Should be considered in conjunction with action and changes

# Example

- *Ana*: Do you want to do something tonight?
- *Bob*: Sure, what do you want to do?
- *Ana*: Let us have a pot-luck dinner with *Christina*. I will prepare some sandwiches and call *Christina*. But can you pick her up? Also, could you bring some soft-drinks?
- *Bob*: Sure. How about 7pm?
- *Ana*: Great.

# Commitments



- Ana commits to prepare sandwich and call Christina.
- Bob commits to get soft-drink, picks up Christina, and both go to Ana's house.
- Reasoning about commitments strongly related to agents' behaviors and capabilities, time, deadlines.

# Questions

- What does it mean for an agent
  - to satisfy a commitment?
  - to violate a commitment?
- How to specify commitment in a declarative manner?
- How to reason about commitment?

# Goal and Contributions

- Investigate the use of action languages in formalizing commitments
- Contributions
  - An action language  $\mathcal{L}^{\text{mt}}$  suitable for reasoning about commitment
  - Uniform view of reasoning about satisfaction, violation, achievement of commitments as reasoning about narratives and planning

# Action Language $\mathcal{L}^{\text{mt}}$

- Domain description
  - over set of agents  $\mathcal{AG}$ , sets of fluents  $\mathcal{F}_i$ , sets of actions  $\mathcal{A}_i$
  - set of statements encoding effects of actions
  - effects of actions: delay, override by latter actions
  - action precondition

# $\mathcal{L}^{\text{mt}}$ Example - Netbill Protocol

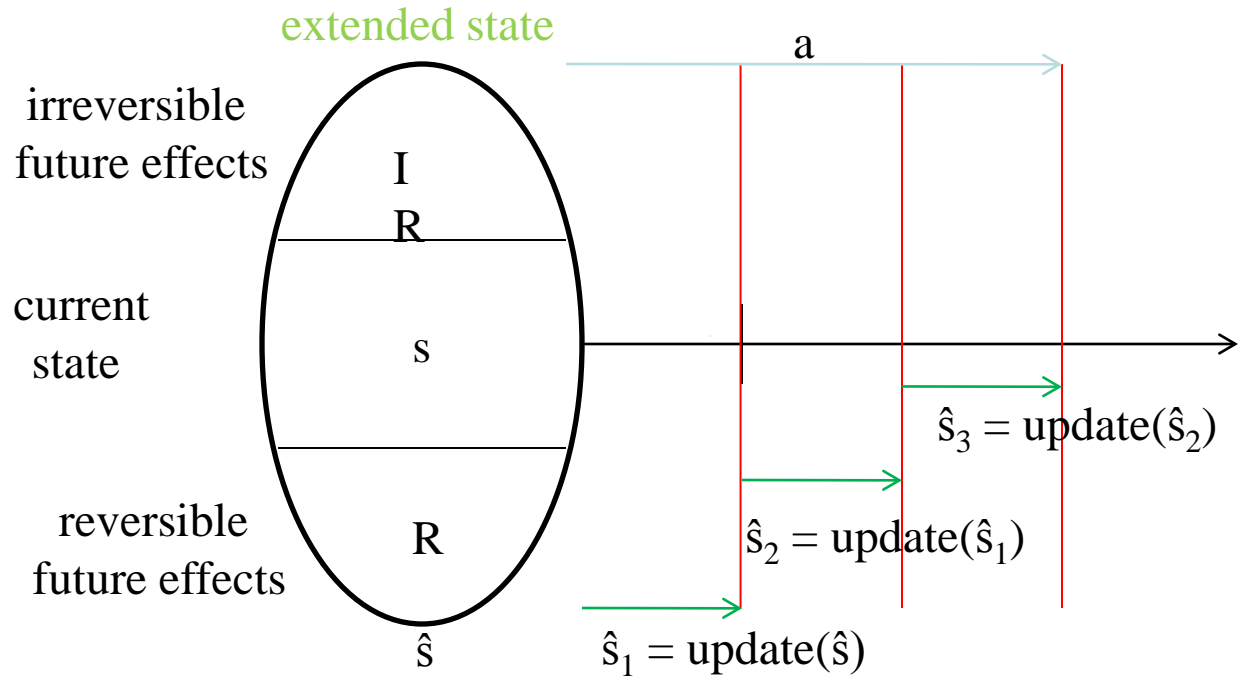
- Customer
  - request for a quote
  - accept a quote if quote has been received
  - send a payment if quote is accepted
  - cancel payment if check has not been deposited
- Merchant
  - send a quote
  - send the goods
  - send the receipt if payment is made



# $\mathcal{L}^{\text{mt}}$ Syntax

- Customer
  - **sendPayment starts** `payment_process` **reversible** `paid`<sup>∨[3,5]</sup>  
[sendPayment starts a reversible process that causes *paid* to be true within 3 to 5 unit of time]
  - **cancelPayment stops** `payment_process`  
[cancelPayment stops the payment process]
  - **impossible** {cancelPayment} **if** `paid`
  - ...
- Merchant
  - **sendQuote starts** `quote_process` **irreversible** `quote`<sup>∨[2,2]</sup>  
[sendQuote starts an irreversible process that causes `quote` to be true after 2 unit of time]
  - ...

# $\mathcal{L}^{\text{mt}}$ Semantics transition function based approach



- $\text{update}(\hat{s})$ : the state after one unit of time
- $\Phi(\hat{s}, \alpha)$ : set of trajectories if  $\alpha$  is executed in  $\hat{s}$
- $\Phi(\hat{s}, \alpha, t) + t_1$ : set of extended states  $t_1$  unit from the time  $\alpha$  is executed which starts after  $t$  unit from  $\hat{s}$

# Representing Observations

$\varphi$ at $s$	$\neg$ paid at $s_0$
$\alpha$ between $s_1, s_2$	$\neg$ accept at $s_0$
$\alpha$ occurs_at $s$	$\neg$ quote at $s_0$
$s_1 < s_2$	$\neg$ good at $s_0$
$s$ at time	sendRequet occurs_at $s_0$
	sendAccept occurs_at $s_1$
	sendGood between $s_2, s_3$
	$s_2$ at 2
	$s_0 < s_1 < s_2 < s_3$

# $\mathcal{L}^{\text{mt}}$ Narratives

- Narrative: (D,O)
  - D: domain specification
  - O: observation
- Specifies a set of possible trajectories

$\neg$ paid at  $s_0$   
 $\neg$ accept at  $s_0$   
 $\neg$ quote at  $s_0$   
 $\neg$ good at  $s_0$

sendRequet occurs\_at  $s_0$

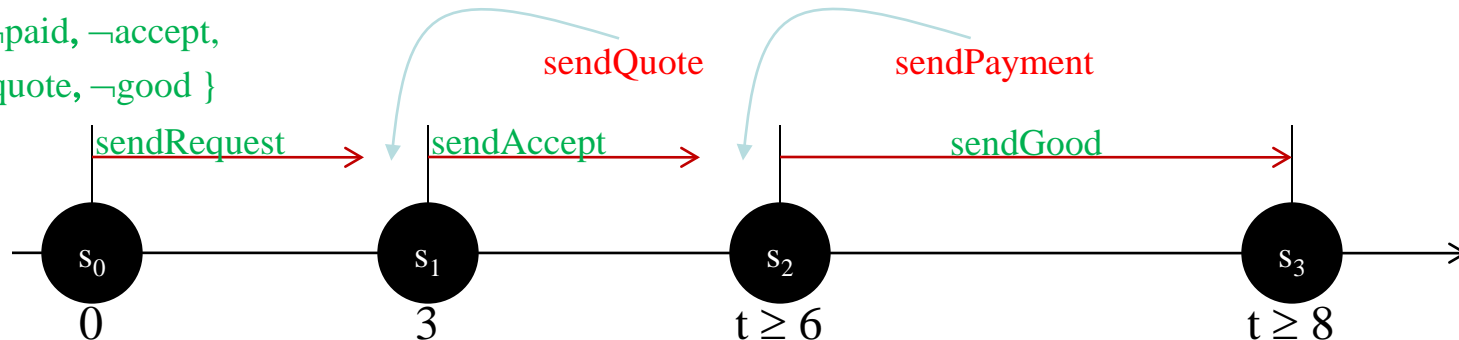
sendAccept occurs\_at  $s_1$

sendGood between  $s_2, s_3$

$s_1$  at 3

$s_0 < s_1 < s_2 < s_3$

$\{ \neg$ paid,  $\neg$ accept,  
 $\neg$ quote,  $\neg$ good  $\}$



# Commitments in $\mathcal{L}^{\text{mt}}$

- Commitment:  $c(x,y,\varphi,t_1,t_2)$ :  $x,y \in \mathcal{AG}$ ,  $\varphi$  - fluent formula,  $t_1 \leq t_2$  - time points
  - $c(\text{Anna}, \text{Bob}, \text{sandwich}, 0, 7)$
  - $c(\text{Bob}, \text{Anna}, \text{arrived}, 0, 7)$
- Activities:
  - $\text{create}(x,y,\varphi,t_1,t_2)$  [x commits to archive  $\varphi$  for y between  $t_1$  and  $t_2$ ]
  - $\text{discharge}(x,y,\varphi)$  [x discharges the commitment  $\varphi$  towards y]
  - $\text{release}(x,y,\varphi)$  [x releases y from the commitment  $\varphi$ ]
  - $\text{assign}(x,y,k,\varphi,t_1,t_2)$  [y transfers  $\varphi$  to another creditor]
  - $\text{delegate}(x,y,k,\varphi,t_1,t_2)$  [x delegates  $\varphi$  to another debtor]
  - $\text{cancel}(x,y,\varphi,\psi,t_1,t_2)$  [x cancels  $\varphi$  and generate a new  $\psi$ ]

# Manipulation of Commitments

- Enabling statements

[ $\varphi|a$ ] **triggers** commitment\_activity

- Examples:

- Merchant agrees to send the receipt between 1 and 3 working days after receiving payment

paid **triggers** c(m,c,receipt,1,3)

- Customer agrees to pay between 5 working days after sending the acceptance of the quote

sendAccept **triggers** create(c,m,paid,1,5)

# Domain with Commitments

- $M = (D, C)$  –  $D$  is a domain description and  $C$  is a set of trigger statements
- Semantics of  $M = (D, C)$  is specified by translating  $M$  to a  $\mathcal{L}^{\text{mt}}$  domain  $M^*$ 
  - a **triggers**  $\text{create}(x, y, \varphi, t_1, t_2)$  is translated into
    - a **causes**  $c(x, y, \varphi)$
    - a **starts**  $c(x, y, \varphi)$  **reversible**  $\text{done}(x, y, \varphi)^{\vee[t_1, t_2]}$   
where  $\text{done}(x, y, \varphi)$  means  $c(x, y, \varphi)$  needs to be realized.
  - ...

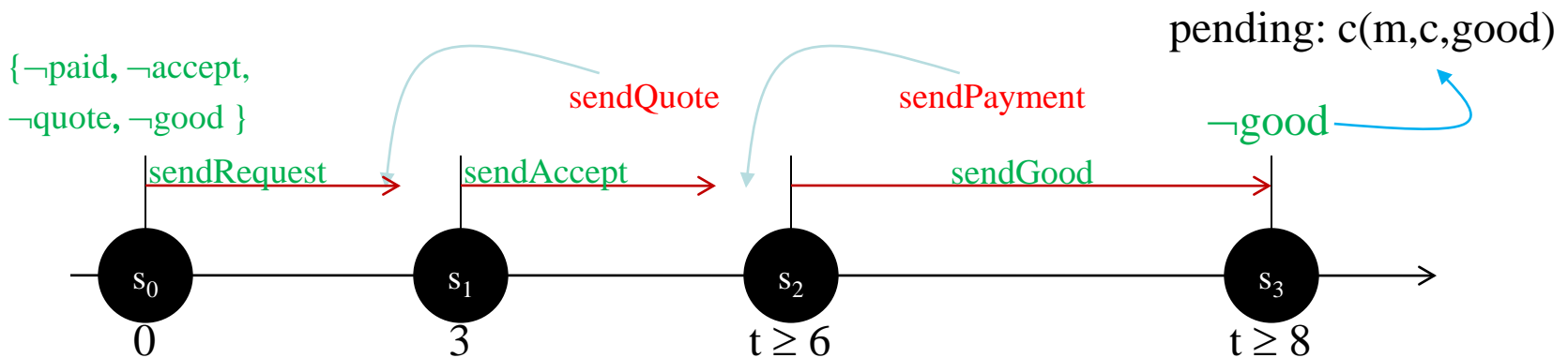
# Commitments w.r.t a Sequence of States

- $M = (D, C)$  and  $[u_0, \dots, u_n]$  is a sequence of states
- $c(x, y, \varphi)$  is
  - satisfied in  $u_n$  if  $u_n \models \neg c(x, y, \varphi)$
  - violated in  $u_n$  if  $u_n \models c(x, y, \varphi) \wedge \text{done}(x, y, \varphi)$
  - pending in  $u_n$  if  $u_n \models c(x, y, \varphi)$  and  
 $u_n \not\models \text{done}(x, y, \varphi)$



# Commitments wrt. a Narrative

- $(D, O, C) - (D, C)$  is a domain with commitments and  $O$  is a set of observations
- $(D, O, C) -$  specifies a set of trajectories  $\rightarrow$  enable the reasoning about satisfied, violated, or pending commitments
- E.g., `sendGood` triggers `create(m,c,good,3,7)`



# Planning for Pending Commitments

- Given  $M = (D, O, C)$
- Identification of pending commitments:  
 $M \models c(x, y, \varphi)$  and  $M \not\models \text{done}(x, y, \varphi)$
- How to satisfy  $c(x, y, \varphi)$ ?  
Planning to achieve  $\neg c(x, y, \varphi)$

# Related Work

- $\mathcal{L}^{\text{mt}}$  first action language with
  - static causal laws [Action language B, Gelfond & Lifschitz]
  - time and deadlines [ADC, Baral et al.]
  - observations [L, Baral et al.]
  - multi-agent domains
  - reversible and irreversible effects
- Commitments
  - basic commitment ontology as in [Mallaya & Huhns, Singh]
  - protocols different from [Giordano et al., Yolum & Singh]

# Conclusions

- Introduced an action language for multi-agent domain with time, deadline, reversible and irreversible effects, and observations.
- Achieved reasoning about commitments and protocols through entailment relation between narratives and queries.
- Our next goal is applying multiagent planning techniques in generating plans to satisfy pending commitments.